# Extremal Functions of Forbidden Matrices PRIMES Conference 

Meghal Gupta<br>Mentor: Jesse Geneson

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## $0,1 \ldots k$-matrix

## Definition

$0,1 \ldots k$-matrix: A matrix where all the entries are in $\{0,1 \ldots k\}$.

## Example

$$
\left(\begin{array}{lll}
1 & 0 & 1 \\
1 & 0 & 1 \\
0 & 0 & 1
\end{array}\right)
$$

is a 0,1 -matrix. It is also a $0,1 \ldots k$-matrix for all $k>1$

## Containment and Avoidance for 0,1-Matrices

$$
\begin{array}{cc}
\mathrm{A} & \mathrm{~B} \\
\left(\begin{array}{lll}
1 & 0 & 1 \\
1 & 0 & 1 \\
0 & 0 & 1
\end{array}\right) & \left(\begin{array}{ll}
1 & 1 \\
1 & 0
\end{array}\right)
\end{array}
$$

## Definition

$A$ contains $B$ : We can delete rows and columns in $A$, and replace 1's with 0's to end up with $B$.

## Containment and Avoidance for 0,1-Matrices

$$
\begin{gathered}
\text { A } \\
\left(\begin{array}{lll}
1 & 0 & 1 \\
1 & 0 & 1 \\
0 & 0 & 1
\end{array}\right) \\
\rightarrow
\end{gathered} \underset{\text { Delete }}{\rightarrow\left(\begin{array}{ll}
1 & 1 \\
1 & 1
\end{array}\right)} \rightarrow \underset{\text { Replace }}{\rightarrow\left(\begin{array}{ll}
1 & 1 \\
1 & 0
\end{array}\right)}
$$

Therefore, $A$ contains $B$.

## Containment and Avoidance for 0,1-Matrices

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0 & 0 & 1
\end{array}\right) \\
\rightarrow\left(\begin{array}{ll}
1 & 1 \\
1 & 1
\end{array}\right) \\
\rightarrow
\end{gathered} \rightarrow \underset{\text { Delete }}{\rightarrow\left(\begin{array}{ll}
1 & 1 \\
1 & 0
\end{array}\right)}
$$

Therefore, $A$ contains $B$.

## Definition

$A$ avoids $B$ : $A$ does not contain $B$.

## Extremal Function 0, 1-matrices

## Definition

ex $(P, n)$ : If $P$ is a 0,1 -matrix, then this is the maximum number of 1 's we can have in an $n \times n$ matrix that avoids $P$.

## Example

$$
P=\left(\begin{array}{ll}
1 & 1 \\
1 & 1
\end{array}\right)
$$

ex $(P, 3)=6$, because any $3 \times 3$ matrix with over 6 ones contains $P$.

## Simple Result

## Theorem

$$
\text { Let } P=\left(\begin{array}{ccccc}
1 & 1 & 1 & . & . \\
0 & 0 & 0 & & \\
0 & 0 & 0 & & \\
. & & & . & \\
. & & & & .
\end{array}\right)
$$

where $P$ has $r$ rows and $c$ columns.
Then, ex $(P, n)=n(r+c)-(r-1)(c-1)$ Important: this is $O(n)$

## Containment and Avoidance for $0,1 \ldots k$-Matrices

$$
\begin{array}{cc}
\mathrm{A} & \mathrm{~B} \\
\left(\begin{array}{lll}
2 & 0 & 1 \\
3 & 1 & 1 \\
0 & 0 & 2
\end{array}\right) & \left(\begin{array}{cc}
2 & 1 \\
0 & 2
\end{array}\right)
\end{array}
$$

## Definition

$A$ contains $B$ : We can delete rows and columns in A, and replace numbers with smaller numbers to end up with $B$.

## Containment and Avoidance for $0,1 \ldots k$-Matrices

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\mathrm{A} & \mathrm{~B} \\
\left(\begin{array}{lll}
2 & 0 & 1 \\
3 & 1 & 1 \\
0 & 0 & 2
\end{array}\right) & \left(\begin{array}{cc}
2 & 1 \\
0 & 2
\end{array}\right)
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## Definition

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## Definition

$A$ avoids $B$ : $A$ does not contain $B$.

## Extremal Function $0,1 \ldots k$-matrices

## Definition

$\mathrm{ex}_{k}(P, n)$ : The maximum sum of numbers an $n \times n 0,1 \ldots k$-matrix $A$ can have and still avoid $P$, where $P$ is a $0,1 \ldots k$-matrix.

## Example

$$
P=\left(\begin{array}{ll}
2 & 1 \\
1 & 1
\end{array}\right)
$$

$\mathrm{ex}_{2}(P, 3)=14$, because any $3 \times 3$ matrix with all entries $0,1,2$ and sum over 14 contains $P$.

## Mapping 0, 1...k-matrices to 0,1-matrices

## Definition

$P_{j}$ : The 0,1 -matrix formed by mapping all entries with values $\geq j$ to 1 and entries $\leq j-1$ to 0 .

## Example

$$
\begin{gathered}
P \\
\left(\begin{array}{lll}
2 & 0 & 1 \\
3 & 1 & 1 \\
0 & 0 & 2
\end{array}\right) \rightarrow\left(\begin{array}{lll}
1 & 0 & 0 \\
1 & 0 & 0 \\
0 & 0 & 1
\end{array}\right)
\end{gathered}
$$

## $0, j$-matrices

## Lemma

Let $P$ be a matrix with only $0, j$ entries. For any $k \geq j$, $e x_{k}(P, n)=(j-1) n^{2}+(k-j+1) e x\left(P_{j}, n\right)$.

## $0, j$-matrices

## Lemma

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## Reasoning.

The 'optimal' matrix that avoids $P$ has $j-1$ entries where an 'optimal' matrix that avoids $P_{j}$ has 0's, and $k$ entries where it has 1's. For example, an 'optimal' matrix that avoids some pattern $P$ might look like:

$$
\left(\begin{array}{ccc}
j-1 & k & k \\
k & k & j-1 \\
k & j-1 & j-1
\end{array}\right)
$$

Calculating, this has $(j-1) n^{2}+(k-j+1) \operatorname{ex}\left(P_{j}, n\right)$ sum.

## Simple Inequality

## Theorem

$(j-1) n^{2}+(k-j+1) e x\left(P_{j}, n\right) \leq \operatorname{ex}_{k}(P, n) \leq$ $(j-1) n^{2}+(k-j+1) e x\left(P_{1}, n\right)$
where $j$ is the maximum element in $P$.

## Simple Inequality

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where $j$ is the maximum element in $P$.

## Proof of LHS.

We find a matrix contained by $P$ that has extremal function $(j-1) n^{2}+(k-j+1) \operatorname{ex}\left(P_{j}, n\right)$.

Let $P^{\prime}$ be the $P$ with all non- $j$ entries replaced with 0 's.

$$
\left(\begin{array}{ccc}
j & 0 & 1 \\
j & j-1 & 1 \\
0 & 0 & j
\end{array}\right) \rightarrow\left(\begin{array}{lll}
j & 0 & 0 \\
j & 0 & 0 \\
0 & 0 & j
\end{array}\right)
$$

## Simple Inequality

## Theorem

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where $j$ is the maximum element in $P$.

## Proof of RHS.

We find a matrix that contains $P$ that has extremal function $(j-1) n^{2}+(k-j+1) \operatorname{ex}\left(P_{1}, n\right)$.

Let $P^{\prime}$ be the $P$ with all non-0 entries replaced with 1's.

$$
\left(\begin{array}{ccc}
j & 0 & 1 \\
j & j-1 & 1 \\
0 & 0 & j
\end{array}\right) \rightarrow\left(\begin{array}{lll}
j & 0 & j \\
j & j & j \\
0 & 0 & j
\end{array}\right)
$$

## Some 0,1,2-matrices

$$
n^{2}+\operatorname{ex}\left(P_{2}, n\right) \leq \operatorname{ex}_{2}(P, n) \leq n^{2}+\operatorname{ex}\left(P_{1}, n\right)
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$$

## Example

$$
\text { Let } P=\left(\begin{array}{ll}
2 & 1 \\
1 & 1
\end{array}\right)
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$\mathrm{ex}_{2}(P, n)=n^{2}+2 n-1 \leftarrow$ analogous to lower bound

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## Example

$$
\text { Let } P=\left(\begin{array}{ll}
2 & 2 \\
2 & 1
\end{array}\right)
$$

$\mathrm{ex}_{2}(P, n) \geq n^{2}+3 n-2 \leftarrow$ NOT analogous to lower bound
We generally believe that the lower bound is closer than the upper bound though.

## A general form

## Theorem

$$
\text { Let } P=\left(\begin{array}{cccccc}
2 & 2 & 2 & . & . & . \\
1 & 1 & 1 & & & \\
1 & 1 & 1 & & & \\
\cdot & & & . & & \\
\cdot & & & & \cdot & \\
\cdot & & & & & .
\end{array}\right)
$$

$e x_{2}(P, n)=n^{2}+O(n)$.

## A general form

## Theorem

$$
\text { Let } P=\left(\begin{array}{cccccc}
2 & 2 & 2 & . & . & . \\
1 & 1 & 1 & & & \\
1 & 1 & 1 & & & \\
\cdot & & & . & & \\
\cdot & & & & \cdot & \\
\cdot & & & & & .
\end{array}\right)
$$

$e x_{2}(P, n)=n^{2}+O(n)$.
Again serves as evidence that lower bound is better:

$$
\begin{gathered}
n^{2}+\operatorname{ex}\left(P_{2}, n\right) \leq \operatorname{ex}_{2}(P, n) \leq n^{2}+\operatorname{ex}\left(P_{1}, n\right) \\
n^{2}+O(n) \leq \operatorname{ex}(P, n) \leq n^{2}+O\left(n^{2-\epsilon}\right)
\end{gathered}
$$

## Improving the Simple Inequality

## Theorem

Let $P$ be a $0,1,2$-matrix. The sum of numbers in a $n \times n$ matrix avoiding $P$ is at most $\leq n^{2}+O\left(k \operatorname{ex}\left(P_{2}, \frac{n}{\sqrt{k}}\right)\right)$, where $k$ is the number of 0 's in the $n \times n$ matrix.
Can be modified to $n^{2}+O\left(\sqrt{k} e x\left(P_{2}, n\right)\right)$, easier to use but a bit weaker.

## Improving the Simple Inequality

## Proof Sketch.

Consider an $n \times n 0,1,2$-matrix that avoids $P$. Build 'boxes' around the 0 's as such:

We can limit the number and side length of the boxes to get the desired result.

## Implications

## Theorem

Let $P$ be a $0,1,2$-matrix. The sum of numbers in a $n \times n$ matrix avoiding $P$ is at most $\leq n^{2}+O\left(k \operatorname{ex}\left(P_{2}, \frac{n}{\sqrt{k}}\right)\right)$, where $k$ is the number of 0 's in the $n \times n$ matrix.
Can be modified to $n^{2}+O\left(\sqrt{k} e x\left(P_{2}, n\right)\right)$, easier to use but a bit weaker.

- Bounds $\mathrm{ex}_{2}(P, n)$ in terms of number of 0 's in the $n \times n$ matrix that must avoid $P: \operatorname{ex}_{2}(P, n) \leq n^{2} O\left(\sqrt{k} \operatorname{ex}\left(P_{2}, n\right)\right)$
- When $k$ is (nontrvially) small, $O\left(\sqrt{k} \operatorname{ex}\left(P_{2}, n\right)\right)$ is better than ex $\left(P_{1}, n\right)$
- When $k$ is large, it is not good enough.
- To make the result useful, we need to find another way to deal with the lots of $k$ 's case.


## Further Research

## Overarching Goal

Characterize all the extremal functions ex $(P, n)$ in terms of 0,1 extremal functions ex $(P, n)$.

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Characterize all the extremal functions ex $x_{k}(P, n)$ in terms of 0,1 extremal functions ex $(P, n)$.

- For 0,1,2-matrices, find $\mathrm{ex}_{2}(P, n)-n^{2}$ upto a constant. We believe it is $\theta\left(\operatorname{ex}\left(\left(P_{2}, n\right)\right)\right.$ rather than $\theta\left(P_{1}, n\right)$


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Characterize all the extremal functions ex $(P, n)$ in terms of 0,1 extremal functions ex $(P, n)$.

- For 0,1,2-matrices, find $\mathrm{ex}_{2}(P, n)-n^{2}$ upto a constant. We believe it is $\theta\left(\operatorname{ex}\left(\left(P_{2}, n\right)\right)\right.$ rather than $\theta\left(P_{1}, n\right)$
- Find the exact value of $\operatorname{ex}_{2}(P, n)$ where

$$
\text { Let } P=\left(\begin{array}{ll}
2 & 2 \\
2 & 1
\end{array}\right)
$$

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Characterize all the extremal functions ex $(P, n)$ in terms of 0,1 extremal functions ex $(P, n)$.

- Generalize the theorem about a row of 2's followed by rows of 1 's to $0,1 \ldots k$-matrices, where we consider a row of $i$ 's followed by rows of $j$ 's where $i>j$.


## Further Research

## Overarching Goal

Characterize all the extremal functions $e x_{k}(P, n)$ in terms of 0,1 extremal functions ex $(P, n)$.

- Generalize the theorem about a row of 2's followed by rows of 1 's to $0,1 \ldots k$-matrices, where we consider a row of $i$ 's followed by rows of $j$ 's where $i>j$.
- Generalize the last theorem (improvement on upper bound from simple inequality) to $0,1 \ldots k$-matrices.


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